Regret of Narendra Shapiro Bandit Algorithms

S. Gadat

Toulouse School of Economics Joint work with F. Panloup and S. Saadane.

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I - Introduction

- I 1 Motivations
- I 2 Stochastic multi-armed bandit model
- I 3 Regret of Stochastic multi-armed bandit algorithms
- I 4 Roadmap

II Narendra Schapiro algorithm (NSa)

- II 1 An historical algorithm (1969)
- II 2 Improvement through penalization
- II 3 Over-penalized NSa

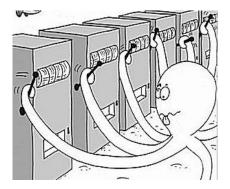
III Weak limit of the Over-penalized NSa

- III 1 Rescaling
- III 2 Trajectories of the rescaled over-penalized NSa
- III 3 Ergodicity and Invariant measure
- III 4 Ergodicity and mixing rate

IV Conclusion

I - 1 Motivations - Stochastic Bandit Games

Problem : You want to earn as much as possible in casino



- You are in a casino and want to play with slot machines
- Each one can give you a potential gain, but these gains are not equivalent
- · You sequentially play with one of the arms of the bandit machine

How to design a good policy to sequentially optimize the gain?

I - 1 Motivations - Dynamic Ressource Allocation

Problem : Optimization of a sequence of clinical trials



Imagine you are a doctor :

- A sequence of patients visit you sequentially (one after another) for a given disease
- You choose one treatment/drug among (say) 5 availables
- The treatments are not equivalent
- You do not know where is the best drug, but you observe the effect of the prescribed treatment on each patient
- You expect to find the best drug despite some uncertainty on the effect of each treatment

How can we design a good sequence of clinical trials?

I - 1 Motivations - Dynamic Ressource Allocation

Problem : "Fast fashion" retailer



Source : Farias & Madan, Operation Research, Vol. 9, No 2, 2011 Imagine you are a firm solding clothes :

- A population of customers visit you sequentially (one after another) each week/day
- You observe weekly/daily sales and measure item's popularity
- · You want to restock popular items and weed out unpopular ones on-line
- You expect to maximize your benefit while finding the best items

How can we design a good sequence of fast-fashion operations?

I - 1 Motivations - Dynamic Ressource Allocation

Other motivating examples

- Pricing a product with uncertain demand to maximize revenue
- Trading (sequentially allocate a ratio of fund to the more efficient trader)
- Recommender systems :
 - advertisement
 - website optimization
 - news, blog posts



Computer experiments

- A code can be simulated in order to optimize a criterion
- This simulation depends on a set of parameters
- Simulation is costly and only few choices of parameters are possible

I - 2 Stochastic multi-armed bandit model Environment :

- At your disposal : d arms with unknown parameters $\theta_1, \ldots, \theta_d$.
- For any time t, and for any choice $I_t \in \{1..., d\}$, you receive a reward :

A_t^{It}

For any choice of an arm *i*, rewards are i.i.d. $(A_t^i)_{t \ge 0} \sim \nu_{\theta_i}$.

Reward distribution :

- In general, the reward distributions ν_{θ} belong to a parametric family (Exponential, Poisson, ...)
- In this talk, simplest case of Bernoulli rewards $\nu_p = \mathcal{B}(p)$:
 - ${}^{\scriptstyle \mathbf{\flat}}$ you obtain a gain of 1 with probability p
 - 0 otherwise (with probability 1 p).
 - Unknown probability of success : (p_1, \ldots, p_d) . Without I.o.g., we assume that

$$p_1 > \max_{2 \leqslant j \leqslant d} p_j.$$

Admissible policy :

The agent's action follow a dynamical strategy, which is defined on-line :

$$I_t = \pi \left(A_{t-1}^{I_{t-1}} \dots, A_1^{I_1} \right).$$

Final goal : Maximize (in expectation) the cumulative rewards :

$$\mathbb{E}\left[\sum_{t=1}^n A_t^{I_t}\right].$$

I - 3 Regret of Stochastic multi-armed bandit algorithms

Regret of an algorithm It yields the minimization of the expected regret R_n

$$\mathbb{E}R_n = \mathbb{E}\max_{1 \le j \le d} \sum_{t=1}^n A_t^j - \mathbb{E}\sum_{t=1}^n A_t^{I_t} = \mathbb{E}\max_{1 \le j \le d} \sum_{t=1}^n (A_t^j - A_t^{I_t})$$

The expectation of the maximum makes the regret difficult to handle, but...

Proposition (Pseudo-regret) If we define $\bar{R}_n := \max_{1 \leq j \leq d} \mathbb{E} \left[\sum_{t=1}^n (A_t^j - A_t^{I_t}) \right]$, one has $\mathbb{E}R_n \leq \bar{R}_n + \sqrt{\frac{n \log d}{2}}.$

This upper bound is useful since

Proposition (Lower bound - (Auer, Cesa-Bianchi, Freund, Schapire 2002)) Uniformly among all policies π and among all Bernoulli distribution rewards :

$$\min_{\pi} \left\{ \max_{\substack{\sup \\ 2 \leqslant j \leqslant d}} \mathbb{E}R_n \right\} \geqslant \frac{\sqrt{nd}}{20}.$$

Conclusion : Upper bounds of \bar{R}_n of the order \sqrt{nd} are competitive (optimal).

I - 4 Roadmap

In this talk, we will :

Briefly describe a standard old-fashioned method

$$X_{t+1} = X_t + \gamma_{t+1}h(X_t) + \gamma_{t+1}\Delta M_{t+1}$$

Introduce a new one whose regret will be studied :

 $\forall n \in \mathbb{N}^* \qquad \bar{R}_n \leq C\sqrt{n}?$

· Provide an asymptotic limit of this penalized bandit up to a correct scaling

$$\beta_n(X_n - \delta_1) \xrightarrow[n \to +\infty]{w^*} \mu$$

 Describe ergodic properties of the rescaled process (Piecewise Deterministic Markov Process)

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- 1 Motivations
- 2 Stochastic multi-armed bandit model
- I 3 Regret of Stochastic multi-armed bandit algorithms
- l 4 Roadmap

II Narendra Schapiro algorithm (NSa)

- II 1 An historical algorithm (1969)
- II 2 Improvement through penalization
- II 3 Over-penalized NSa

III Weak limit of the Over-penalized NSa

- III 1 Rescaling
- III 2 Trajectories of the rescaled over-penalized NSa
- III 3 Ergodicity and Invariant measure
- III 4 Ergodicity and mixing rate

IV Conclusion

II - 1 An historical algorithm (1969)

The so-called Narendra-Shapiro bandit algorithm (NSa for short) defines a probability vector of \mathcal{S}_d

$$X_t = (X_t^1, \dots, X_t^d) \qquad | \qquad \sum_{j=1}^a X_t^j = 1.$$

Idea : Use X_t to sample one arm at step t and then upgrade this probability X_t according to the obtained reward to deduce X_{t+1} .

• In the two-armed situation with $p_2 < p_1$, denote $X_t = (x_t, 1 - x_t)$

$$x_{t+1} = x_t + \begin{cases} \gamma_{t+1}(1-x_t) & \text{if player } 1 \text{ is selected and wins} \\ -\gamma_{t+1}x_t & \text{if player } 2 \text{ is selected and wins} \\ 0 & \text{otherwise} \end{cases}$$

• Multi-armed situation, I_t : arm sampled at time t, $A_t^{I_t}$: obtained reward. Upgrade

$$\forall j \in \{1 \dots d\} \qquad X_t^j = X_{t-1}^j + \gamma_t \left[\mathbf{1}_{\{I_t = j\}} - X_{t-1}^j \right] A_t^{I_t}$$

- To sum up :
 - If you win : reinforce the probability to sample I_t w.r.t. the remaining weights $(X_t^j)_{j \neq I_t}$ and decrease the probability to sample the other arms accordingly.
 - If you loose $(A_t^{I_t} = 0)$: do nothing.
- Common step size :

$$\gamma_t = (1 + t/C)^{-\alpha}$$
, $\alpha \in (0, 1)$ with large enough C.

II - 1 An historical algorithm (1969)

 ${\sf Few words \ about \ NSa:}$

- Recursive stochastic algorithm
- Anytime policy
- Involves nontrivial mathematical difficulties

It can be written as mean drift + martingale increment

$$X_{t+1} = X_t + \gamma_{t+1} h(X_t) + \gamma_{t+1} \Delta M_{t+1}.$$

 x_3

 e_1

x2

 e_2

 x_1

In the 2-armed setting $(p_2 < p_1 \text{ and } X_t = (x_t, 1 - x_t))$, the drift on x_t is

$$h(x) = (p_1 - p_2)x(1 - x).$$

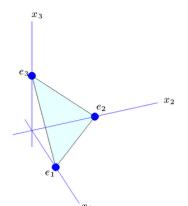
II - 1 An historical algorithm (1969)

$$X_t = (x_t, 1 - x_t)$$

$$x_{t+1} = x_t + \gamma_{t+1} h(x_t) + \gamma_{t+1} \Delta M_t.$$

with

$$h(x) = (p_1 - p_2)x(1 - x)$$



- O.D.E. approximation $\dot{x} = h(x)$, local trap at $\{0\}$ and stable equilibrium at $\{1\}$.
- But : the conditional variance term vanishes at 0 and 1, making impossible the use of Duflo's argument about the escape of local traps.
- $\succ \text{ Indeed, for any sequence } \gamma_t = \left(\frac{C}{t+C}\right)^{\alpha}, \qquad \alpha \in (0,1), \text{ the algorithm is fallible}$

$$\mathbb{P}\left(\lim x_t = 0\right) > 0 \Longrightarrow \mathbb{E}R_n \gtrsim Cn \gg \sqrt{n}$$

II - 2 Improvement through penalization

What's wrong with NSa? Gittins, JRSS(B)'79 :

Good regret properties only occur with an exploration/exploitation trade-off...

- NSa is almost a pure exploitation method : no exploration term to exit local traps.
- Main idea : Introduce a penalty term [Lamberton & Pages, EJP'09]
- In the 2-armed settings $(p_2 < p_1 \text{ and } X_t = (x_t, 1 x_t))$:

$$X_{t+1} = X_t + \begin{cases} +\gamma_{t+1}(1-X_t) & \text{if arm 1 is selected and wins} \\ -\gamma_{t+1}X_t & \text{if arm 2 is selected and wins} \\ -\rho_{t+1}\gamma_{t+1}X_t & \text{if arm 1 is selected and loses} \\ +\rho_{t+1}\gamma_{t+1}(1-X_t) & \text{if arm 2 is selected and loses} \end{cases}$$



When one arm fails, decrease the probability to sample it.

LP'09 : Up to technical conditions on (ρ_t, γ_t) : penalized 2-armed bandit is infallible (a.s. convergence to the good target)

II - 3 Over-penalized NSa

This additional penalty term will be inefficient from the minimax regret point of view. As a last resort : increase the penalty effect to reinforce the escape from local traps :

$$X_{t+1} = X_t + \begin{cases} +\gamma_{t+1}(1-X_t) - \rho_{t+1}\gamma_{t+1}X_t & \text{if arm 1 is selected and wins} \\ -\gamma_{t+1}X_t + \rho_{t+1}\gamma_{t+1}(1-X_t) & \text{if arm 2 is selected and wins} \\ -\rho_{t+1}\gamma_{t+1}X_t & \text{if arm 1 is selected and loses} \\ +\rho_{t+1}\gamma_{t+1}(1-X_t) & \text{if arm 2 is selected and loses} \end{cases}$$

Whatever happens with the selected arm, it is penalized (escape from local traps).



A multi-armed version :

$$\begin{aligned} X_{t}^{j} &= X_{t-1}^{j} + \gamma_{t} \left[\mathbf{1}_{I_{t}=j} - X_{t-1}^{j} \right] A_{t}^{I_{t}} \\ &- \gamma_{t} \rho_{t} X_{t-1}^{I_{t}} \left[\mathbf{1}_{I_{t}=j} - \frac{1 - \mathbf{1}_{I_{t}=j}}{d - 1} \right] \end{aligned}$$

II - 3 Over-penalized NSa and infallibility Write $X_t = X_{t-1} + \gamma_t h(X_t) + \gamma_t \rho_t \kappa(X_t) + \gamma_t \Delta M_t$. Drift :

$$h^i(x_1,\ldots,x_d)=x_i\left[(1-x_i)p_i-\sum_{j
eq i}x_jp_j
ight],orall i\in\{1,\ldots,d\}$$

Equilibria of $\dot{X} = h(X)$: Dirac masses on each arm. Stable one : (1, 0, ..., 0). The Kushner-Clarck theorem \implies a.s. convergence towards an equilibrium (which ?) Theorem (Infallibility of the Over-penalized NSa) If $p_d \leq p_{d-1} \leq ... \leq p_2 < p_1$ and $\gamma_t = \gamma_1 t^{-\alpha}$, $\rho_t = \rho_1 t^{-\beta}$, then $0 \leq \beta \leq \alpha$ and $\alpha + \beta \leq 1 \Longrightarrow \lim_{t \to +\infty} X_t = (1, 0 ..., 0)$ a.s.

Sketch of proof : The penalty term induced by κ is

$$\kappa^{i}(x) = -x_{i}^{2}(1-p_{i}) + \frac{1}{d-1}\sum_{j \neq i} x_{j}^{2}(1-p_{j}), \forall i \in \{1, \dots, d\}$$

If $X^1_\infty=0,\,\kappa^1(X_\infty)>0$ and :

$$\alpha \leqslant \beta \Longrightarrow \limsup \frac{\sum_t \gamma_t \Delta M_t}{\sum \gamma_t \rho_t} \ge 0$$

$$\alpha + \beta \leqslant 1 \Longrightarrow \sum \gamma_t \rho_t = +\infty \Longrightarrow \sum \gamma_t \rho_t \kappa(X_t) = +\infty$$

II - 4 Non-asymptotic upper bound of the regret

We detail the picture for the two-armed over-penalized NSa

$$\bar{R}_n = \max_{j \in \{1,2\}} \mathbb{E} \sum_{t=1}^n A_t^j - A_t^{I_t}$$
$$= \mathbb{E} \sum_{t=1}^n \left[p_1 - (X_t^1 p_1 + (1 - X_t^1) p_2) \right]$$
$$= (p_1 - p_2) \sum_{t=1}^n \rho_t \underbrace{\frac{1 - X_t^1}{\rho_t}}_{:=Y_t}$$

In S.A., we expect a "CLT" $Y_n \xrightarrow[n \to +\infty]{w^*} \mu$ and if $\sup_{n \ge 1} \mathbb{E} Y_n < \infty$, then

$$\bar{R}_n \lesssim \sum_{t=1}^n \rho_t$$

Choose β in $\rho_t = \frac{\rho_1}{t^{\beta}}$ as large as possible s.t. $\beta \leq \alpha, \alpha + \beta \leq 1$. Optimal calibration :

$$\gamma_t = \frac{\gamma_1}{\sqrt{t}}$$
 and $\rho_t = \frac{\rho_1}{\sqrt{t}}$

II - 4 Non-asymptotic upper bound of the regret

We are turned to the study of the random dynamical system induced by $(Y_t)_{t \ge 1}$. Again :

$$Y_{t+1} = Y_t + \gamma_t \varphi_t(Y_t) + \gamma_t \Delta M_{t+1}.$$

Beyond the analytic formula of φ_t , a simple picture :

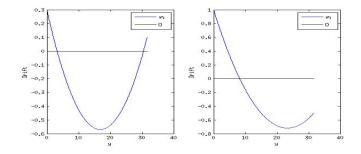


FIGURE: Drift for non penalized (left) and overpenalized (right) NSa when $y \in [0, \gamma_t^{-1}]$.

To control the increments of Y_t , the right situation is much better :

Large value of Y_t are naturally decreased by φ_t

II - 4 Non-asymptotic upper bound of the regret

- Difficulty : obtaining a uniform bound over all the values $0 \le p_2 < p_1 \le 1$.
- · Lyapunov arguments and painful computations lead to non asymptotic bound.
- Key quantity that induces the understanding of the good scaling

 $\pi = p_1 - p_2.$

Theorem (Upper bound of the regret : 2-armed over-penalized NSa)

 $\forall n \in \mathbb{N} \qquad \sup_{p_2 < p_1} \bar{R}_n \leqslant 30\sqrt{2n}.$

Optimal settings : $\gamma_n = \frac{9}{10\sqrt{n}}$ and $\rho_n = \frac{1}{3\sqrt{n}}$. Sketch of proof : Define $Z_t^{(r)} = \frac{(1-X_t)^r}{\gamma_t}$ and exhibit a mean-reverting effect for r sufficiently large $\mathbb{E}[Z_{t+1}^{(r)} | \mathcal{F}_t] = Z_t^{(r)} + P_{t-r}(Z_t^{(r)}).$

Find r such that $P_{t,r}$ is negative on $[C(\gamma_t, \pi), \gamma_t^{-1}]$ where $C(\gamma_t, \pi) = o(\gamma_t^{-1})$ and

$$\sup_{t \ge 0} \mathbb{E}[Z_t^{(r)}] < \infty.$$

• Exhibit a recursion between $\mathbb{E}[Z_t^{(r)}]$ and $\mathbb{E}[Z_t^{(r-1)}]$ for a result on $\sup_{t \ge 0} \mathbb{E}[Y_t]$

II - 4 Numerical simulations

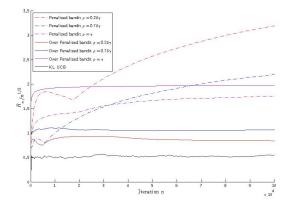


FIGURE: Evolution of $n \mapsto \sup_{(p_1, p_2) \in [0, 1], p_2 \leq p_1} \frac{\bar{R}_n}{\sqrt{n}}$ for over-penalized NSa (continuous colored line) and penalized NSa (dashed colored line) and KL UCB (black line).

- Over-penalization is important for a competitive regret
- Practical : $\bar{R}_n \leqslant \sqrt{n}$ Theoretical : $\bar{R}_n \leqslant 30\sqrt{2n}$
- \triangleright Defeated by UCB-like algorithms from a statistical regret point of view $(\bar{R}_n \leqslant \sqrt{n}/2)$
- Computation time very low compared to MOSS or UCB-like algorithms

II - 4 Numerical simulations

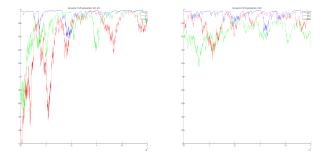


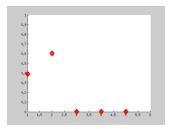
FIGURE: Evolution of the probability of Arm 1 (best one) with respect to n while $p_1 - p_2 = 0.1$. Left : ρ_1/γ_1 is varying. Right : p_2 is increasing.

Seems to behave quite particularly (maybe after a good rescaling?)

- Some jumps randomly distributed ? (more or less frequent according to the parameters)
- ${\scriptstyle \rm P}$ Almost deterministic evolution between jumps when n is large
- Much more faster than KL-UCB (ratio of time needed : 1/100).

II - 4 Numerical simulations

Time for a short movie ... 5 arms, p = [0.9, 0.88, 0.8, 0.75, 0.7].



Let's go back to the mathematics

I - Introduction

- I 1 Motivations
- I 2 Stochastic multi-armed bandit model
- l 3 Regret of Stochastic multi-armed bandit algorithms
- l 4 Roadmap

II Narendra Schapiro algorithm (NSa)

- II 1 An historical algorithm (1969)
- II 2 Improvement through penalization
- II 3 Over-penalized NSa

III Weak limit of the Over-penalized NSa

- III 1 Rescaling
- III 2 Trajectories of the rescaled over-penalized NSa
- III 3 Ergodicity and Invariant measure
- III 4 Ergodicity and mixing rate

IV Conclusion

III - 1 Rescaling

We fix $p_1 > \max(p_2, \ldots, p_d)$, the "good" rescaling of what is left over by X_n^1 is

$$\bar{X}_n = \frac{(X_n^2, \dots, X_n^d)}{\rho_n}$$

Proposition

For any $f \in \mathcal{C}^2(\mathbb{R}^{d-1},\mathbb{R})$:

$$\mathbb{E}\left[f(\bar{X}_{n+1})|\mathcal{F}_n\right] = f(\bar{X}_n) + \gamma_{n+1}\mathcal{L}_d(f)(\bar{X}_n) + o_P(\gamma_{n+1}),$$

where \mathcal{L}_d is the Markov generator given by

$$\mathcal{L}_{d}(f)(\bar{x}) = \sum_{j=2}^{d} \underbrace{\frac{p_{j}}{g} \bar{x}_{j}}_{jump \text{ rate}} \underbrace{\left[f(\bar{x} + g\mathbf{1}_{j}) - f(\bar{x})\right]}_{jump \text{ size}} + \sum_{j=2}^{d} \underbrace{\left[\frac{1 - p_{1}}{d - 1} - p_{1}\bar{x}_{j}\right]}_{deterministic \text{ part}} \partial_{j}f(\bar{x}) \cdot \frac{1 - p_{1}\bar{x}_{j}}{deterministic \text{ part}}$$

- The amount of jump is low when $g = \frac{\gamma_1}{\rho_1}$ is large (seen in simulations).
- ▶ The size of jumps is large when g is large.

III - 1 Rescaling

As a tensorized process, it is enough to study the following Markov generator :

 $\mathcal{L}(f)(\bar{x}) = (a - b\bar{x})f'(\bar{x}) + cx[f(\bar{x} + g) - f(\bar{x})]$

- Family of Piecewise Deterministic Markov Process (PDMP for short)
- Random dynamical systems with an increasing interest (encountered in many modelisation problems)
- Famous examples (among many others) :
 - Telegraph process [Kac, '74]
 - Storage models [Roberts & Tweedie,'00]
 - Randomly switched ODE [Benaïm et al.,'14] & Parrondo-like paradox
 - TCP models [Guillin, Malrieu et al.'13, Cloez & Hairer'13]

What the dynamic looks like exactly in the over-penalized NSa case?

Set

$$a = \frac{1 - p_1}{d - 1}, b = p_1, c_j = \frac{p_j}{g}, g = \frac{\gamma_1}{\rho_1}$$

- Between jumps, the evolution is deterministic and follow a differential flow

$$\dot{\phi}(\xi,t) = \left[\frac{1-p_1}{d-1} - p_1\xi\right]\partial_\xi\phi(\xi,t)$$

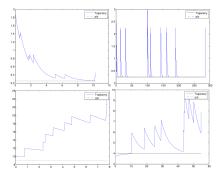
• Poisson jumps with an instantaneous average push of $\frac{p_j}{g}\bar{x}_j \times g$. Here, the size of the jumps are deterministic.

III - 2 Trajectories of the rescaled over-penalized NSa

- \mathcal{L}_d acts as a tensorized Markov generator on each coordinate.
- > The problem is reduced to the study of the random dynamic system described by

$$\mathcal{L}(f)(\bar{x}) = (a - b\bar{x})f'(\bar{x}) + cx[f(\bar{x} + g) - f(\bar{x})],$$

• Examples of rescaled trajectories for several values of (a, b, c, g)



 \triangleright Asymptotic direction : a/b. Bottom left : transient behaviour when $cg > b \ldots$ but in the bandit algorithm

$$cg - b = p_j - p_1 < 0$$
 (!)

III - 3 Ergodicity and Invariant measure

Ergodicity can be helpful to derive confidence bounds. It requires to obtain some mixing properties around an/the invariant measure.

$$\mathcal{L}(f)(\bar{x}) = (a - bx)f'(\bar{x}) + cx[f(\bar{x} + g) - f(\bar{x})],$$

For over-penalized NSa, the process should be studied only when cg - b < 0. Proposition (Invariant measure - rescaled over-penalized NSa) The PDMP \bar{X}_t has a unique invariant measure μ supported by

$$\left[\frac{1-p_1}{p_1(d-1)}, +\infty\right]^{d-1}$$

Sketch of proof : existence and uniqueness through a Lyapunov certificate :

$$\mathcal{L}(Id) = a - (b - cg)Id.$$

But ... Some real difficulties :

- > No explicit formula for μ ... We are far from a standard CLT with a Gaussian distribution and even far from the simplest case of the TCP process
- Less is known about the smoothness of μ ... Intricate situation as pointed by [Bakhtin & Hurth & Mattingly '14].

III - 4 Ergodicity and mixing rate

 \mathcal{L} is a non-reversible Markov operator, which is usual for this kind of kinetic models The question : Obtaining an upper bound of the mixing rate :

$$d(L(X_t),\mu) \leq \epsilon(t) \longrightarrow 0$$
 as $t \longrightarrow +\infty$.

Traditional distance

$$\|L(X_t) - \mu\|_{\mathbb{L}^2(\mu)\mathfrak{O}} = \sup_{f : \|f\|_{\mathbb{L}^2(\mu)} = 1} \|\mathbb{E}[f(\bar{X}_t^x)] - \mu(f)\|_{\mathbb{L}_2(\mu)}$$

Non-reversible generators : difficult to handle with the \mathbb{L}^2 distance, require informations on μ (Modified norms [Villani,'09], Lie brackets [Gadat & Miclo'13])

 Resort less sophisticated distances induced by trajectorial properties (instead of functional ones)
 Wasserstein distance :

$$\mathcal{W}_{p}(\nu_{1},\nu_{2}) = \inf \left\{ \mathbb{E}\left((X-Y)^{p} \right) \right)^{\frac{1}{p}} | L(X) = \nu_{1}, L(Y) = \nu_{2} \right\}$$

Total Variation distance :

$$d_{TV}(\nu_1, \nu_2) = \max_{\Omega \subset E} |\nu_1(\Omega) - \nu_2(\Omega)|$$

Use some coupling techniques to derive quantiative bounds

III - 4 Ergodicity and mixing rate

The simple idea :

- \triangleright Build a non independent coupling (\bar{X}_t,Y_t) such that \bar{X}_t and Y_t follow the dynamic given by $\mathcal L$ and $Y_0\sim \mu$
- Try to make \bar{X}_t and Y_t close to each others for the Wasserstein results

Theorem (Wasserstein ergodicity)

An explicit constant γ_p exists such that

$$\mathcal{W}_p(L(\bar{X}_t),\mu) \leqslant \gamma_p e^{-t\pi/p},$$

where $\pi = p_1 - p_2$ is the difference between the 2 probabilities of success of the 2 best arms

Optimal for W_1 . Open questions for W_p .

• Try to make the two processes $\bar{X}_t = Y_t$ stucked rapidly for the TV results

Theorem (Total Variation ergodicity)

Some explicit constants C and α exist such that

 $d_{TV}(L(\bar{X}_t),\mu) \leq Ce^{-\alpha\pi t}.$

Suspected to be far from the optimal exponents.

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- I 1 Motivations
- I 2 Stochastic multi-armed bandit model
- l 3 Regret of Stochastic multi-armed bandit algorithms
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IV Conclusion

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Statistics :

- Standard NSa Algorithm is fallible
- Penalized bandits are infallible
- Over-penalization : relevant for regret bounds
- Over-penalization : traduces a vanishing repelling effect on each corner of the simplex.
- Minimax result in the two-armed case :

$\bar{R}_n \leqslant C\sqrt{2n},$

 Much more faster than what is already existing in Bandit methods while statistically competitive (not as good as KL UCB)

Probability :

- Rescaled process as a PDMP.
- Random jumps come from the binary rewards given by each arm.
- Ergodic properties

Anecdotal :

• Used in some trading firms in « La Defense » . . .

IV Conclusion

Open questions :

 \blacktriangleright Regret with d arms? Numerical simulations lead to the conjecture

$$\bar{R}_n \leqslant C\sqrt{dn},$$

which is the known minimax rate for d-armed bandit.



Over-Penalized NSa seems to behave well

- What should be a generalization of Over-Penalized NSa for continuous rewards?
 What is the rescaled process (suspected to be a diffusion instead of a jump process ...)
- Many challenging questions with the PDMP :
 - Spectral results and L² convergence
 - Wasserstein lower bounds
 - Smoothness of the invariant measure

Thank you for your attention