

# Conditional quantile sequential estimation for stochastic codes

T. Labopin-Richard   F. Gamboa   A. Garivier

Institut de mathématiques de Toulouse

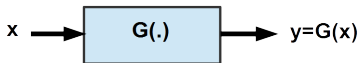
March 23, 2016

# Table of contents

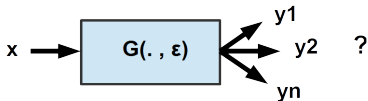
- 1 Problem and overview
  - First ideas
  - To another strategy
- 2 Stochastic algorithms embeded in  $k$ -nearest neighbors method
  - The algorithm
  - Results
- 3 Numerical simulations
  - Dimension 1
  - Dimensions 2 and 3

- 1 Problem and overview
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  - Dimension 1
  - Dimensions 2 and 3

# What is a stochastic code ?

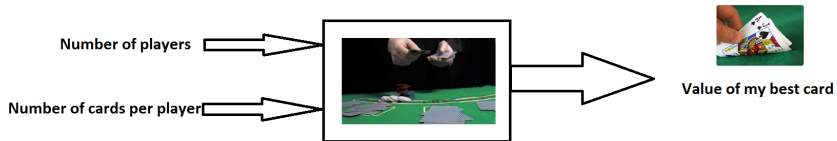


Numerical code :  $Y=G(X)$   
 $G(x)$  is a real number

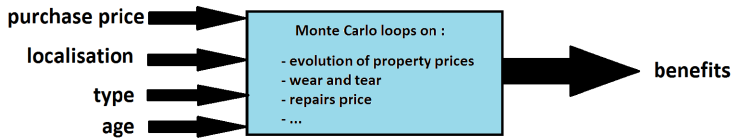


Stochastic code :  $Y=G(X, \epsilon)$   
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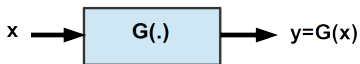
# Example 1 :



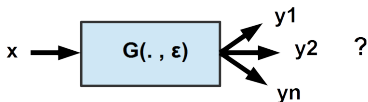
## Example : property investment



# What is a stochastic code ?

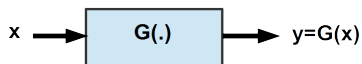


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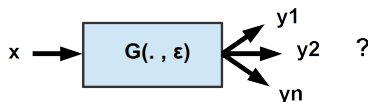


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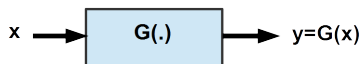


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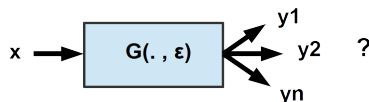
Notation :  $X \in \mathbb{X} \subset \mathbb{R}^d$ .



# What is a stochastic code ?



Numerical code :  $Y=G(X)$   
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Stochastic code :  $Y=G(X, \epsilon)$   
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**Notation :**  $X \in \mathbb{X} \subset \mathbb{R}^d$ .

**Goal :** Estimate the **quantile** of level  $\alpha \in ]0, 1[$  of the law  $\mathcal{L}(G(x, \epsilon))$  using as few as possible calls to the code.

- 1 Problem and overview
  - First ideas
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# To estimate the quantile of a law $Z$ , one can

- 1) Build a sample  $(Z_1, \dots, Z_n)$  of  $Z$ .
- 2) Use a quantile estimator as :
  - a) The empirical quantile  $Z_{(\lfloor n\alpha \rfloor + 1)}$ .  
 $\Rightarrow$  consistant and normaly Gaussian.
  - b) The Robbins-Monro stochastic algorithm

$$\begin{cases} \theta_0 \in \mathbb{R} \\ \theta_{n+1} = \theta_n - \frac{1}{n^\gamma} (\mathbf{1}_{Z_{n+1} \leq \theta_n} - \alpha). \end{cases}$$

$\Rightarrow$  consistant and normaly Gaussian if  $\gamma \in ]1/2, 1]$ .

# Application to our problem

We aim at estimate the quantile of the law  $\mathcal{L}(G(x, \epsilon))$ , we could then

- 1) Provide several times the same input  $x$  to the stochastic code, to build a sample of the target law.
- 2) Use one of the previous estimator.

## Problem

We aim at estimating the conditional quantile for **every input  $x$** .  
If  $\mathbb{X}$  is uncountable or if each call to the code is **expensive**, the is not a solution.

- 1 Problem and overview
  - First ideas
  - To another strategy
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## Type of strategy

- 1) A budget  $N$  of calls to the code is fixed.
- 2) We sample  $(X_1, \dots, X_N)$  from the input law  $X$ .
- 3) We observe the corresponding responses  $(Y_1, \dots, Y_N)$ .
- 4) We apply an algorithm wich allows, for every  $x$  and using only the previous observations, to estimate the conditional quantile.

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  - First ideas
  - To another strategy
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  - The algorithm
  - Results
- 3 Numerical simulations
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## Algorithm

$$\begin{cases} \theta_0 & \in \mathbb{R} \\ \theta_{n+1} & = \theta_n - \frac{1}{n^\gamma} \left( \mathbf{1}_{Z_{n+1} \leq \theta_n} - \alpha \right) \end{cases}$$



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où  $k_n = \lfloor n^\beta \rfloor$ .

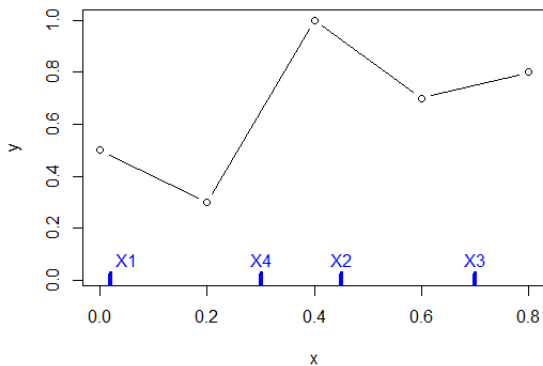
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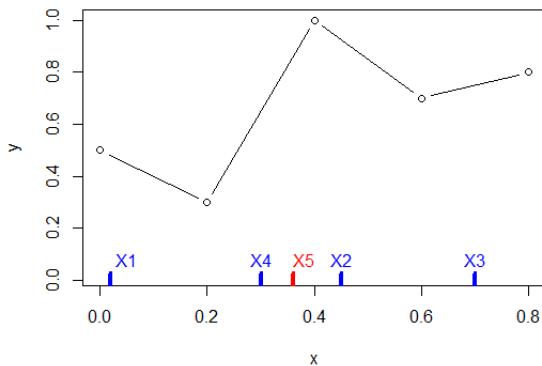
où  $k_n = \lfloor n^\beta \rfloor$ .

- For which parameters  $(\gamma, \beta)$  is the algorithm **convergent**?  
 $\Rightarrow$  Compromise needed.
- Could we prove some non-asymptotic results?
- Are there **optimal parameters**?

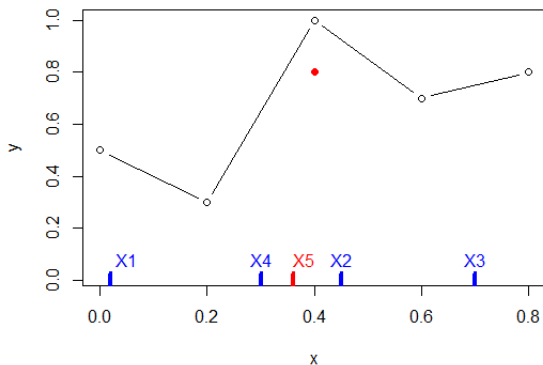
# Example



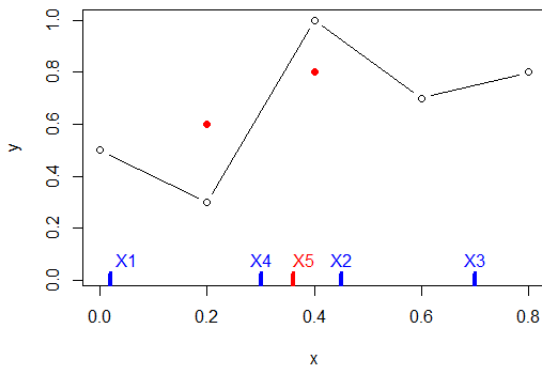
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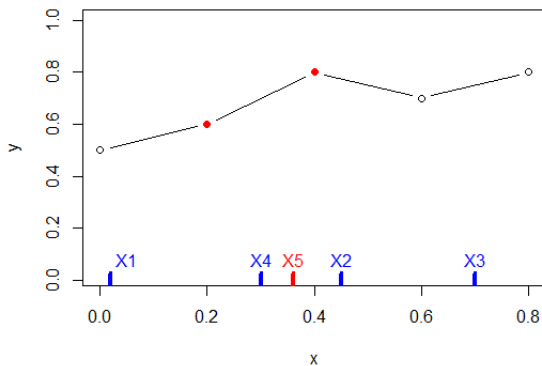


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- 1 Problem and overview
  - First ideas
  - To another strategy
- 2 Stochastic algorithms embeded in  $k$ -nearest neighbors method
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- 3 Numerical simulations
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# Continuity assumption

Notations :

- $\mathcal{B}_x$  the set of the balls of  $\mathbb{R}^d$  centered in  $x$ .  $B \in \mathcal{B}_X$  has a radius  $r_B$ .
- For  $B \in \mathcal{B}_x$ ,  $F_Y^B$  is the cumulative distribution function of the law  $\mathcal{L}(g(X, \epsilon) | X \in B)$ .
- $F_{Y^x}$  is the cumulative distribution function of the law  $\mathcal{L}(g(x, \epsilon))$ .

**Assumption A1** For all  $x \in \text{Supp}(X)$ , there exists a constant  $M(x)$  such that

$$\forall B \in \mathcal{B}_x, \forall t \in \mathbb{R}, |F_{Y^B}(t) - F_{Y^x}(t)| \leq M(x)r_B .$$

## Technical assumptions

**Assumption A2** The input law  $X$  has a density which is lower bounded by a constant  $C_{inputs} > 0$ .

$\Rightarrow$  Useful to deal with  $\mathbb{E}(\|X - x\|_{(k_n, n)})$  or  $\mathbb{P}(X \in kNN_n(x))$ .

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**Assumption A4** For every  $x \in \text{Supp}(X)$ , the law of  $g(x, \epsilon)$  has a density which is lower bounded by a constant  $C_g(x) > 0$ .

$\Rightarrow$  There exists a constant  $C_2(x, \alpha)$  such that :

$$\forall \theta_n(x), [F_{Y^x}(\theta_n(x)) - F_{Y^x}(\theta^*(x))] [\theta_n(x) - \theta^*(x)] \geq C_2(x, \alpha) [\theta_n(x) - \theta^*(x)]^2.$$

# Results

## Theorem : almost-sure convergence

Let  $x$  be a fixed input. Under assumptions **A1** and **A2**, the algorithm in  $x$  is almost surely convergent whenever

$$\frac{1}{2} < \gamma < \beta < 1.$$

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- $\beta < 1 \Rightarrow$  technical assumption.
- $\gamma < \beta \Rightarrow$  *effective learning rate*.

## Classical algorithm :

- An update at each step
- $\sum_n \gamma_n = +\infty \Rightarrow \gamma \in ]\frac{1}{2}, 1]$ .

## Our algorithm :

- An update every  $\frac{k_n}{n} \sim n^{\beta-1}$  steps
- At step  $n$  :  $N = \sum_{k \leq n} k^{\beta-1} \sim n^\beta$  updates
- At time  $t = n^{\frac{1}{\beta}}$  :  $N = n$  updates
- *Effective learning rate* :  $\gamma_{k_n} = \frac{1}{\left(n^{\frac{1}{\beta}}\right)^\gamma}$ .
- $\sum_n \gamma_{k_n} = +\infty \Rightarrow \gamma < \beta$ .

# Rate of convergence of the MSE

## Theorem

Under hypothesis **A1**, **A2**, **A3** and **A4**, the MSE  $a_n(x)$  satisfies :  
 $\forall(\gamma, \beta, \epsilon)$  such that  $0 < \gamma \leq \beta < 1$  and  $1 > \epsilon > 1 - \beta$ ,  $\forall n \geq N_0$ ,

$$a_n(x) \leq \exp\left(-2C_2(x, \alpha) \sum_{j=N_0+1}^n j^{-\epsilon-\gamma}\right) C_1$$

$$+ \sum_{k=N_0+1}^n \exp\left(-2C_2(x, \alpha) \sum_{j=k+1}^n j^{-\epsilon-\gamma}\right) d_k + C_1 \exp\left(-\frac{3n^{1-\epsilon}}{8}\right)$$

where

$$d_n = C_1 \exp\left(-\frac{3n^{1-\epsilon}}{8}\right) + 2\sqrt{C_1} M(x) C_3(d) \gamma_n \left(\frac{k_n}{n}\right)^{\frac{1}{d}+1} + \gamma_n^2 \frac{k_n}{n}.$$

# Compromise between the two errors

- The **bias error** gives the term

$$\exp \left( -2C_2(x, \alpha)(x) \sum_{k=N_0+1}^n \frac{1}{k^{\epsilon+\gamma}} \right).$$

This term decreases to 0 if and only if  $\gamma + \epsilon < 1$  which implies  $\beta > \gamma$ . Then  $\beta$  **has to be chosen not too small**.

- The **online learning error** gives the term

$$\gamma_n (k_n/n)^{1/d+1} = n^{(1-\beta)(1+1/d)+\gamma}.$$

We then need that  $\beta$  is as small as possible compared to 1.  
Then  $\beta$  **has to be chosen not too big**.

## Corollary : optimal parameters

Under the same assumptions than in Theorem 1, the mean square error decreases faster when parameters are  $\gamma = \frac{1}{1+d}$  and  $\beta = \gamma + \eta_\beta$  where  $\eta_\beta > 0$  is as small as possible. Moreover, with these parameters, there exists a constant  $C_9(x, \alpha, d)$  such that  $\forall n \geq N_4(x, \alpha, d)$ ,

$$a_n(x) \leq \frac{C_9(x, \alpha, d)}{n^{\frac{1}{1+d} - \eta}}$$

where  $\eta = \frac{\eta_\epsilon}{2} + \eta_\beta$  and  $\eta_\epsilon = 1 - \beta - \epsilon$ .

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  - To another strategy
- 2 Stochastic algorithms embeded in  $k$ -nearest neighbors method
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# Dimension 1 - almost-sure convergence

Two models with  $X \sim \mathcal{U}([-1, 1])$ ,  $\epsilon \sim \mathcal{U}([-0.5, 0.5])$  and  $x = 0$  :

$$g(X, \epsilon) = X^2 + \epsilon \text{ et } g(X, \epsilon) = |X| + \epsilon$$

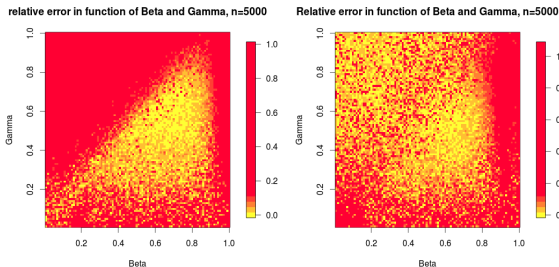


FIGURE – Almost-sure convergence in function of  $\beta$  et  $\gamma$ .

# Mean Square Error

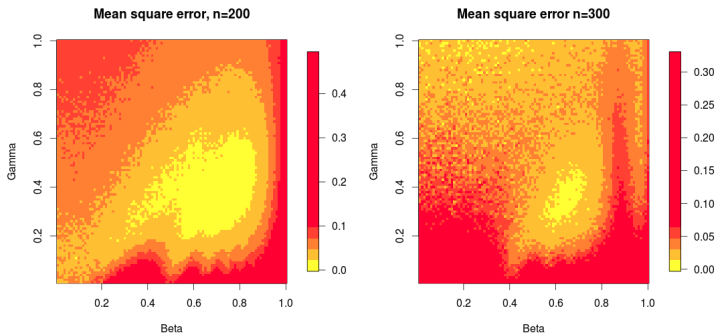


FIGURE – Convergence of the mean square error in function of  $\beta$  et  $\gamma$ .

- 1 Problem and overview
  - First ideas
  - To another strategy
- 2 Stochastic algorithms embeded in  $k$ -nearest neighbors method
  - The algorithm
  - Results
- 3 Numerical simulations
  - Dimension 1
  - Dimensions 2 and 3

## Dimensions 2 and 3

Two models with  $g(X, \epsilon) = \|X\|^2 + \epsilon$  for  $X \sim \mathcal{U}([-1, 1]^d)$ ,  
 $\epsilon \sim \mathcal{U}([-0.5, 0.5])$  and  $x = 0_{\mathbb{R}^d}$ .

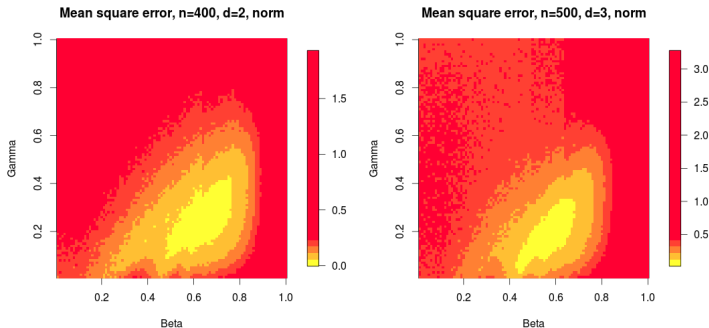


FIGURE – Mean Square Error in function of  $\beta$  and  $\gamma$ .

# Conclusion et perspectives

## Conclusion :

- We introduced an algorithm to estimate the conditional quantile of the output law of a stochastic code.
- We give the best parameters to tune the algorithm (to reach the best rate of convergence of the MSE).
- Numerical simulations show that our algorithm is powerful to solve the problem.

## Perspectives :

- What is happening if we relax the compact support assumptions ?
- Could we find lower bound for the Mean Square Error ?
- Apply this algorithm to real data.

Thanks for your attention.